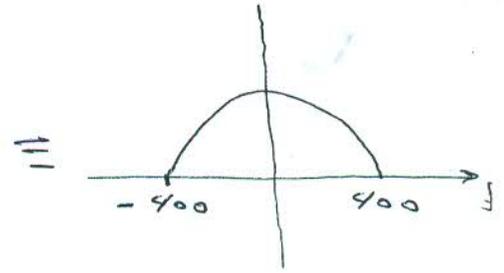


رابطه اتصال (1) نظم الاتصالات (2) داسنامه عمري

[1] Question one (20 mark)

A- (7 marks)

$$M(f) = \begin{cases} \cos\left(\frac{\pi f}{800}\right) & -400 \leq f \leq 400 \\ 0 & \text{o.w} \end{cases}$$

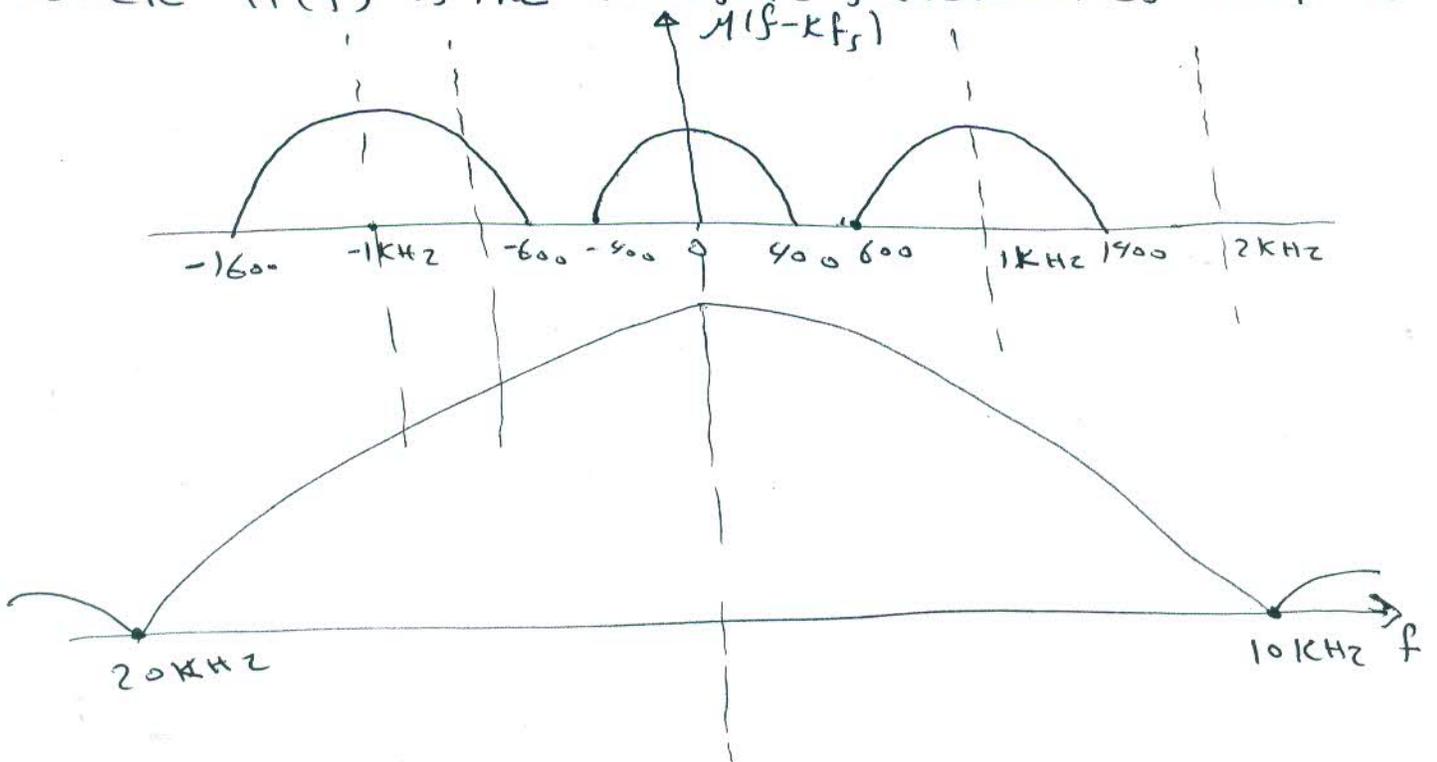


$$f_s = 1 \text{ KHz}, T_{\text{pulse}} = 0.1 \text{ ms}$$

$$M(f) = \cos\left(\frac{\pi f}{800}\right) \cdot \text{rect}\left(\frac{f}{800}\right)$$

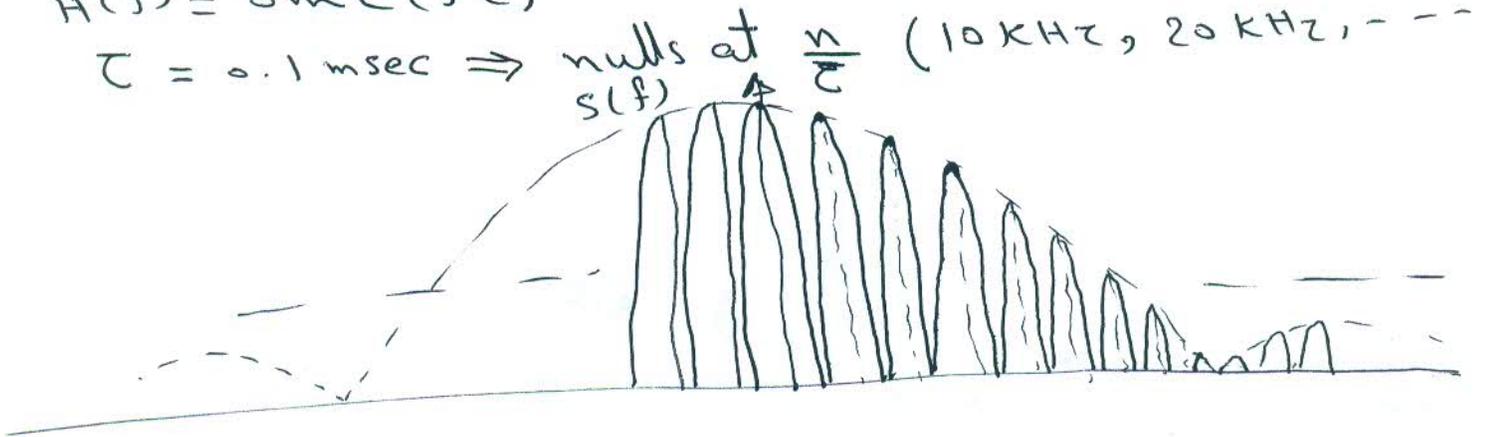
The spectrum of PAM: $S(f) = f_s \sum_{k=-\infty}^{\infty} M(f - k f_s) \cdot H(f)$

where H(f) is the transfer function of the pulse



$$H(f) = \text{sinc}(f\tau)$$

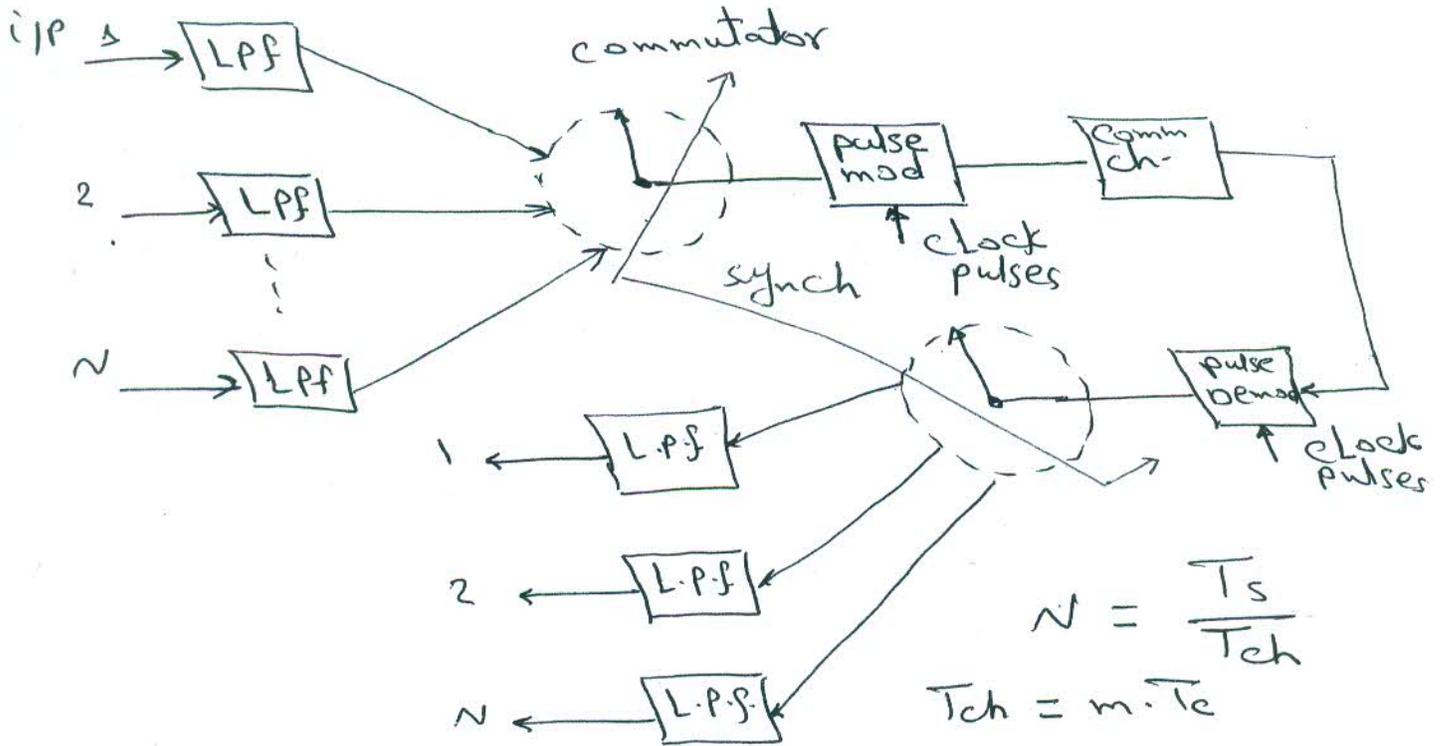
$\tau = 0.1 \text{ msec} \Rightarrow$ nulls at $\frac{n}{\tau}$ (10KHz, 20KHz, ...)



B- (8 marks)

i → The principles of TDM.

- enables the joint utilization of a common ch. by a plurality of independent message sources without mutual interference

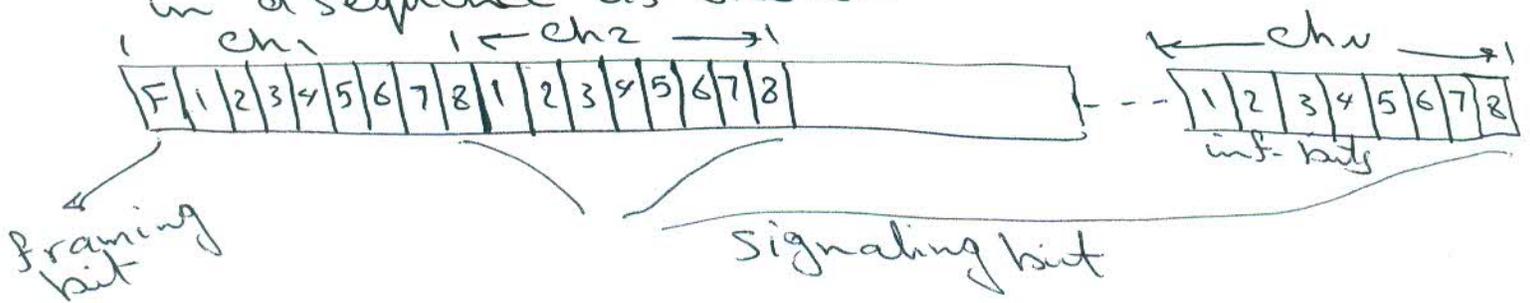


→ pulse mod. transform the multiplexed signal into a form suitable for Txion over the comm. ch.

→ TDM introduces a B.W expansion factor N

→ in application using PCM, each sample represented by code word

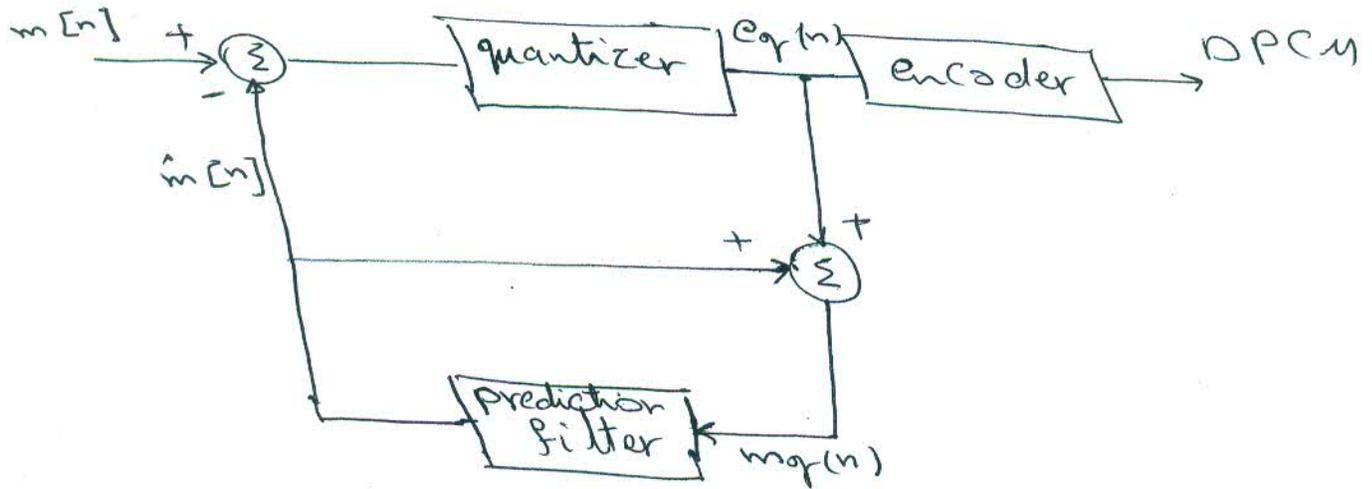
→ Binary code words of each ch. are multiplexed in a sequence as shown



(3)

c - show that the performance of DPCM can be controlled by the design of prediction filter.

for DPCM



$m[n]$: unquantized i/p

$e[n] = m[n] - \hat{m}[n]$ → prediction of $m[n]$

prediction error

$$e_q[n] = e[n] + q[n]$$

⇒ using DPCM results in processing gain

$$(SNR)_0 = \frac{\sigma_M^2}{\sigma_Q^2} \quad \begin{array}{l} \sigma_M^2 : \text{variance of the original i/p} \\ \sigma_Q^2 : \text{variance of } q[n] \end{array}$$

$$(SNR)_0 = \left(\frac{\sigma_M^2}{\sigma_E^2} \right) \cdot \left[\frac{\sigma_E^2}{\sigma_Q^2} \right] = G_p \cdot (SNR)_Q$$

σ_E^2 : variance of the prediction error

$(SNR)_Q$: signal to quantization noise ratio

$$G_p = \frac{\sigma_M^2}{\sigma_E^2} > 1$$

for higher G_p = design the prediction filter to minimize σ_E^2

ii_

$$N = 24$$

TDM

$$T_{\text{pulse}} = 1 \mu\text{s}$$

and extra pulse for synch

$$f_m = 3.4 \text{ KHz}$$

$$\text{Let } f_s = 8 \text{ KHz} > 2 f_m$$

$$\therefore T_s = \frac{1}{f_s} = 125 \mu\text{sec}$$

= spacing

each channel assigned time T_{ch}

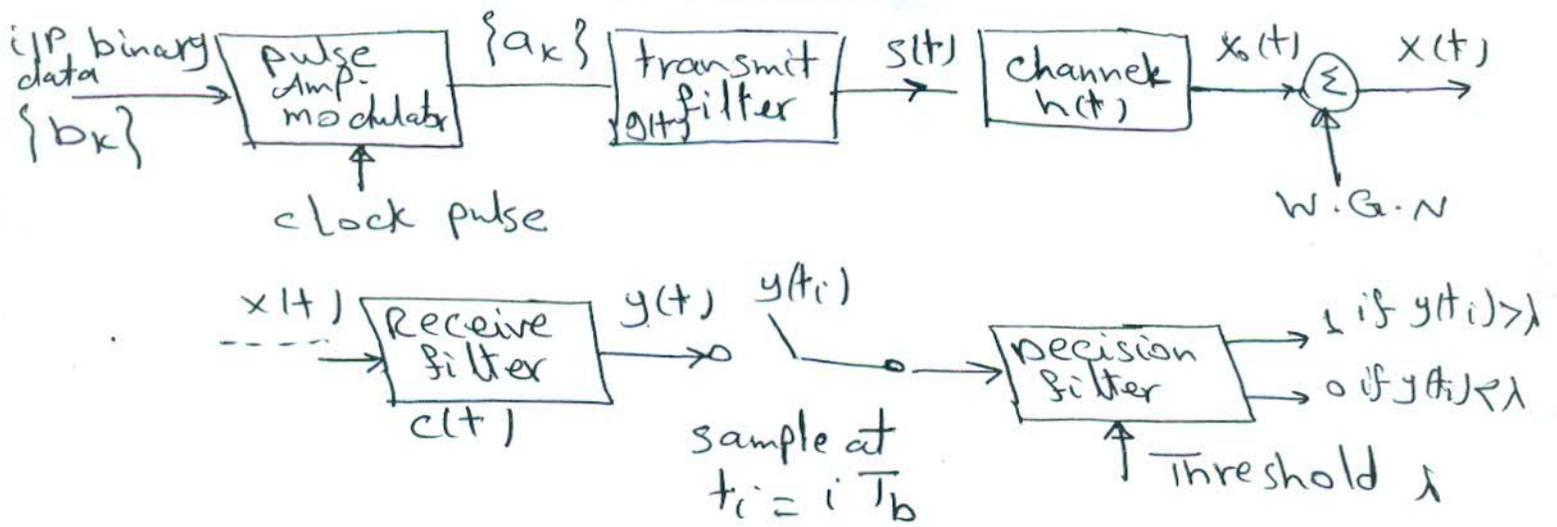
$$T_{\text{ch}} = \frac{T_s}{N + 1 \text{ synch}} = \frac{125}{25} = 5 \mu\text{sec}$$

since the pulse duration = $1 \mu\text{sec}$

$$\therefore \boxed{S = 4 \mu\text{sec}}$$

[2] Question 2 : (20 Mark) (4)

A ^{7 marks} ISI, arises when the comm ch. is dispersive.
 - consider a base band binary PAM system



$$a_k = \begin{cases} +1 & \text{if } b_k \text{ is } 1 \\ -1 & \text{if } b_k \text{ is } 0 \end{cases}$$

$$s(t) = \sum_k a_k g(t - kT_b)$$

→ The receiver filter o/p

$$y(t) = \mu \sum_k a_k p(t - kT_b) + n(t)$$

μ : scaling factor accounts for amplitude changes.

$$\mu p(t) = g(t) * h(t) * c(t), \quad p(0) = 1$$

$$\therefore \mu p(f) = G(f) \cdot H(f) \cdot C(f)$$

$n(t)$: noise produced at the o/p of the receiver filter due to channel noise $w(t)$

→ the o/p $y(t)$ is sampled at $t_i = iT_b$

$$y(t_i) = \mu \sum_{k=-\infty}^{\infty} a_k p[(i-k)T_b] + n(t_i)$$

$$= \mu a_i + \mu \sum_{k=-\infty}^{\infty} a_k p[(i-k)T_b] + n(t_i)$$

→ μa_i represent the contribution of the i th fixed bit
- 2nd term represents the residual effect of all other
fixed bits $\equiv \bar{I} S \bar{I}$

Then to eliminate the effect of $\bar{I} S \bar{I}$

$$P(i T_b - k T_b) = \begin{cases} 1 & i = k \\ 0 & i \neq k \end{cases}$$

B-(6)

(5)

$$NRZ \Rightarrow P_e = 10^{-6}$$

- Signaling rate is doubled $\Rightarrow T_b' = \frac{1}{2} T_b$

$$P_e = \frac{1}{2} \operatorname{erfc} \left(\sqrt{\frac{E_b}{N_0}} \right) = \frac{1}{2} \operatorname{erfc}(u)$$

$$\therefore u = 3.3$$

when the signaling rate is doubled

$$\therefore E_b' = A^2 T_b' = A^2 \cdot \frac{T_b}{2}$$

$$\therefore P_e' = \frac{1}{2} \operatorname{erfc} \left(\frac{u}{\sqrt{2}} \right) = \frac{1}{2} \operatorname{erfc}(2.33)$$

$$= 10^{-3}$$

C-(7 marks)

an equalizer

$$C(\omega) = \left[a_1 e^{-j\omega T} + a_2 e^{-j2\omega T} \right] M(\omega)$$

\therefore The transfer function of the channel

$$H_{ch}(\omega) = a_1 e^{-j\omega T} + a_2 e^{-j2\omega T}$$

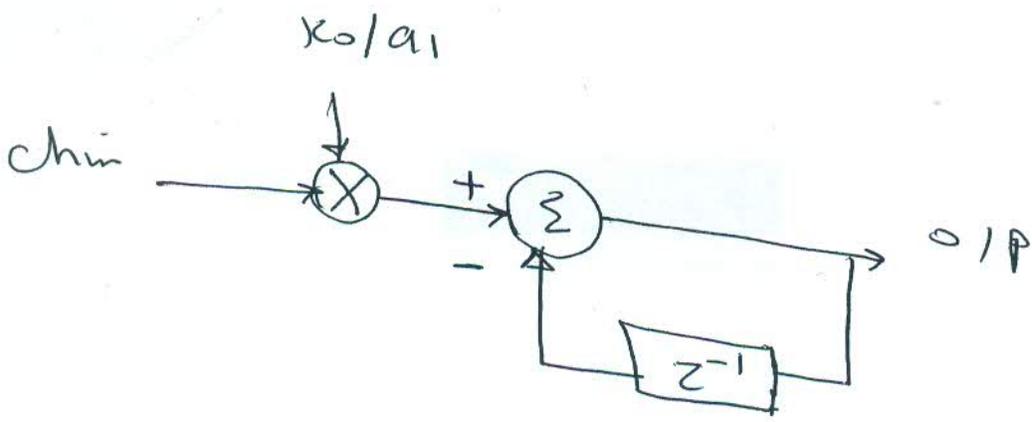
\Rightarrow The equalizer should be designed so that

$$H_{ch}(\omega) \cdot H_e(\omega) = k_0 \exp(-j\omega T)$$

where T is the transmission delay

$$\begin{aligned} \therefore H_e(\omega) &= \frac{k_0 \exp(-j\omega T)}{a_1 e^{-j\omega T} + a_2 e^{-j2\omega T}} \\ &= \frac{(k_0/a_1)}{1 + \frac{a_2}{a_1} e^{-j\omega T}} \end{aligned}$$

\swarrow
O/P
I/P



[Q3]

[6]

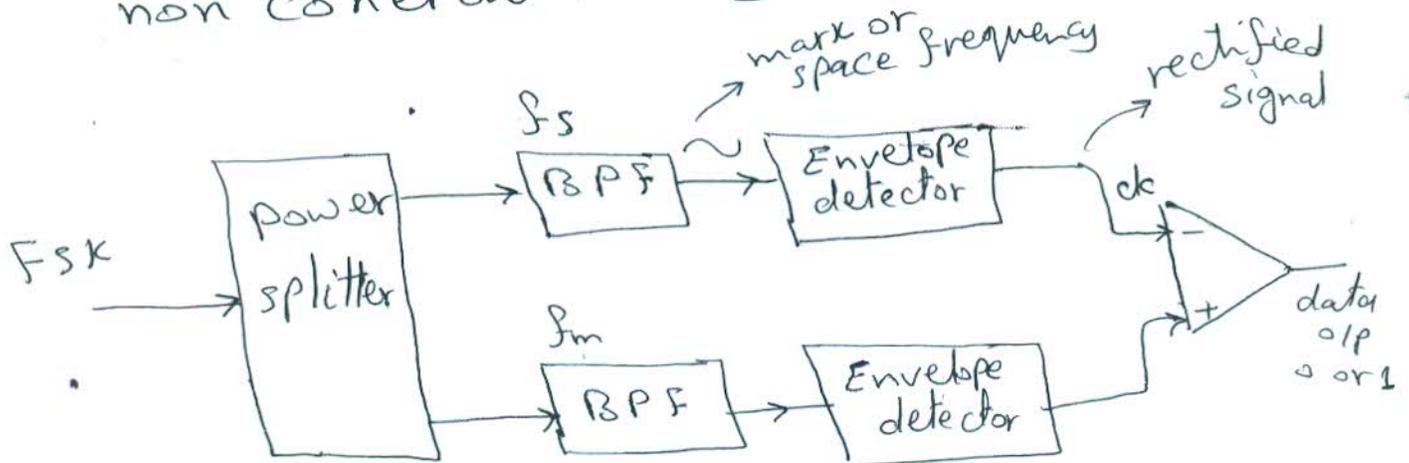
A- 6 marks

⇒ metrics for choice of digital modulation scheme

- 1- High spectral efficiency $\eta_b = \frac{R_b}{B_T}$
- 2- " power efficiency $\eta_p = \frac{E_b}{N_0}$
- 3- " data rates : Bits per second
- 4- Robust to multipath effects, fading conditions
- 5- Easy to implement and cost effective
- 6- Low carrier-to-cochannel signal interference ratio
- 7- Low out of band radiation
- 8- constant or near constant envelope

B- 6 marks

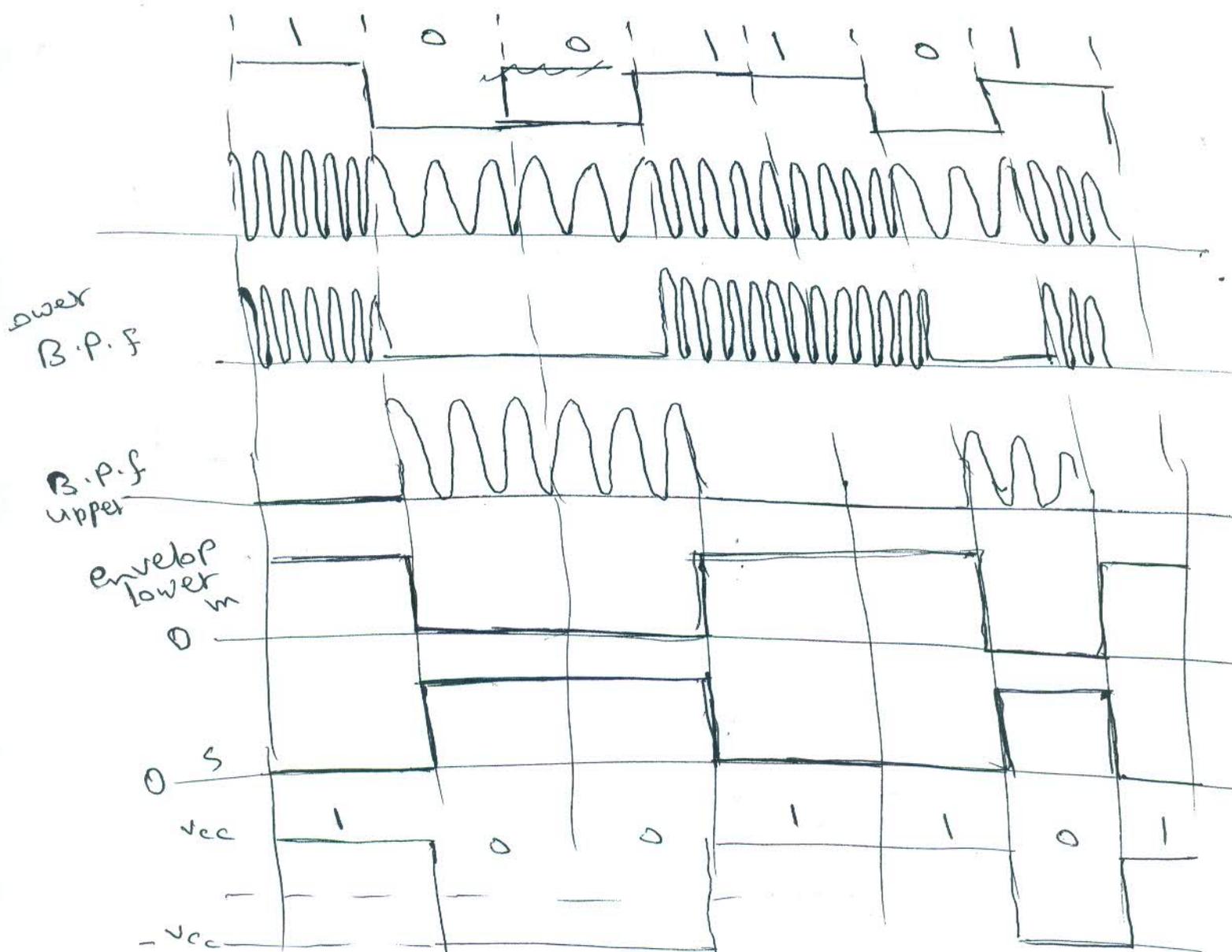
⇒ The operation of the demodulator for non coherent binary FSK.



→ no frequency is involved in the demodulator process that is synchronized either in phase, frequency or both with the incoming FSK

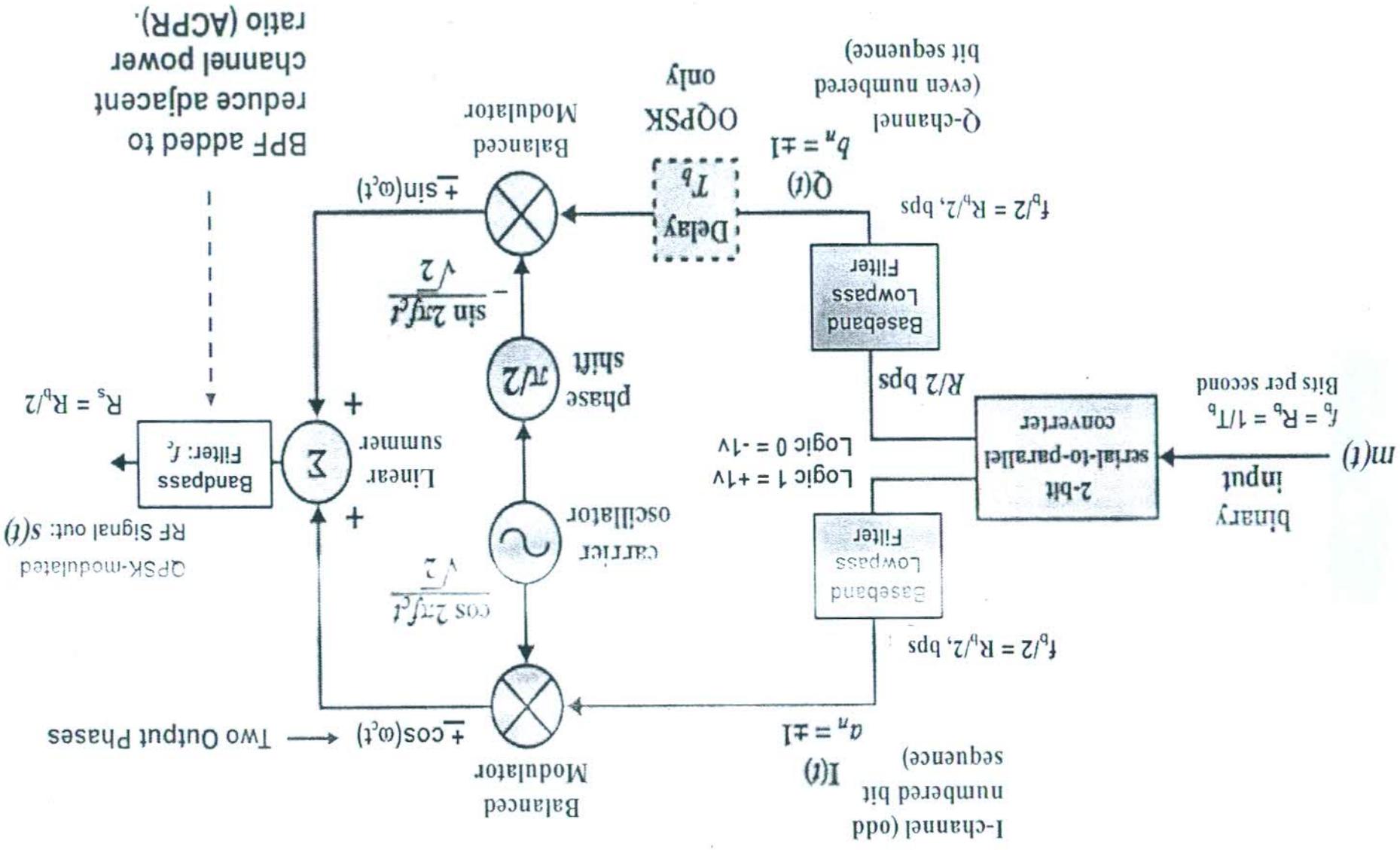
[7]

- no carrier recovery circuits
- The operation the incoming signal is splitted by power splitter
- The upper B.P.F output is related to space freq (0 bits) while the lower is related to mark freq (1 bit)
- The envelop produces dc level corresponds to either mark or space



Transmit Modulator

For Quadrature Phase Shift Keying & Offset-QPSK modulation



(9)

(c)

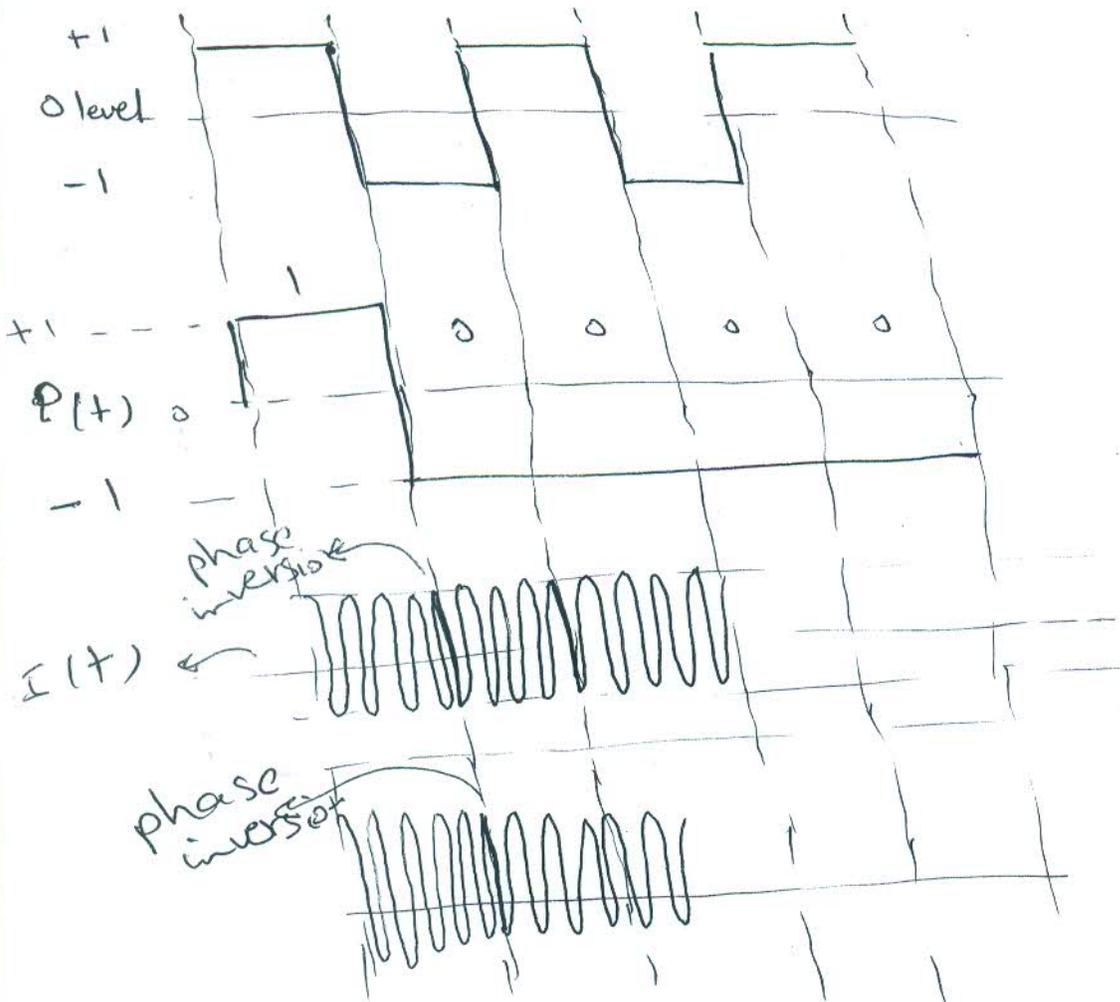
1	1	0	0	1	0	0	0	1	0
↑	↑		↑	↑	↑				

$I(t) \rightarrow 10101$

$Q(t) \rightarrow 10000$

$I(t)$ wave

1 0 1 0 1



for the total ΔIP

the first pair $(I, Q) \equiv (1, 1)$

$$\Delta IP = \frac{\cos \omega_c t}{\sqrt{2}} + \frac{\sin \omega_c t}{\sqrt{2}}$$

angle 45°

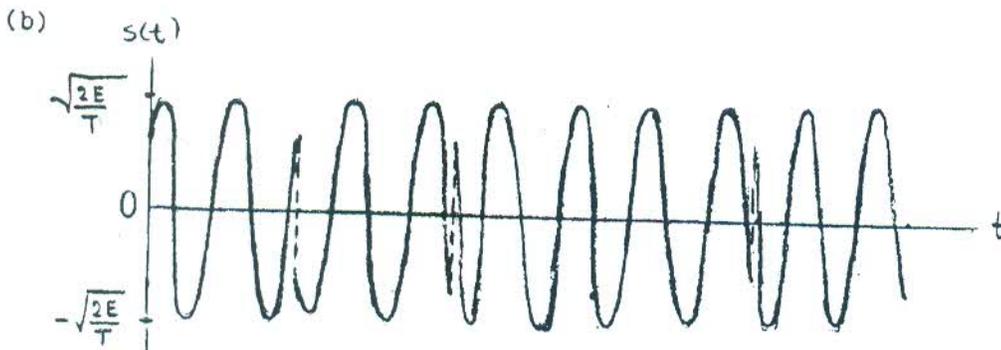
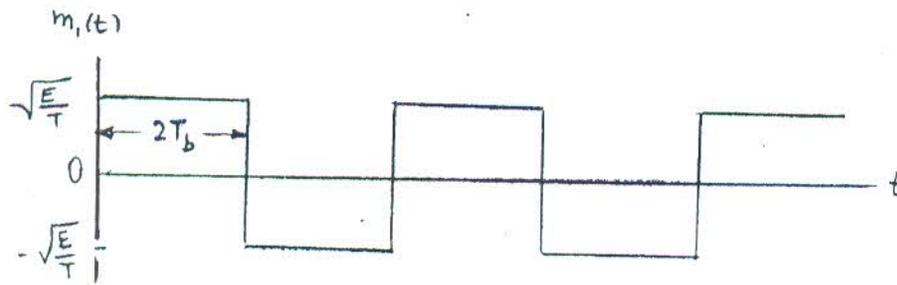
Q3 (c)

(a) The QPSK wave can be expressed as

$$s(t) = m_1(t) \cos(2\pi f_c t) + m_2(t) \sin(2\pi f_c t).$$

Dividing the binary wave into dibits and finding $m_1(t)$ and $m_2(t)$ for each dibit:

dibit	11	00	10	00	10
$m_1(t)$	$\sqrt{E/T}$	$-\sqrt{E/T}$	$\sqrt{E/T}$	$-\sqrt{E/T}$	$\sqrt{E/T}$
$m_2(t)$	$\sqrt{E/T}$	$-\sqrt{E/T}$	$-\sqrt{E/T}$	$-\sqrt{E/T}$	$-\sqrt{E/T}$



[Q4] 20 Marks

- State the main differences between error correction block codes and convolutional codes
- Draw both the encoder shift registers and the encoder state diagram of the convolutional encoder with $(n,k,K)=(2,1,3)$
- The previous encoder in part b is used to encode the following sequence 110010. What is the output coded sequence?
- If an error occurred in the received sequence of part c, explain how the decoder can detect and correct this error.

Solution

A) 6 marks

for block codes,

- On the transmission end, each k -bit block of data is mapped into an n -bit block ($n > k$) called a code word. The first K -bits represents data and the remaining $(n-k)$ represents check bits.
- block code process data in blocks of k bits at a time
- useful when data are transmitted and received in discrete manner.

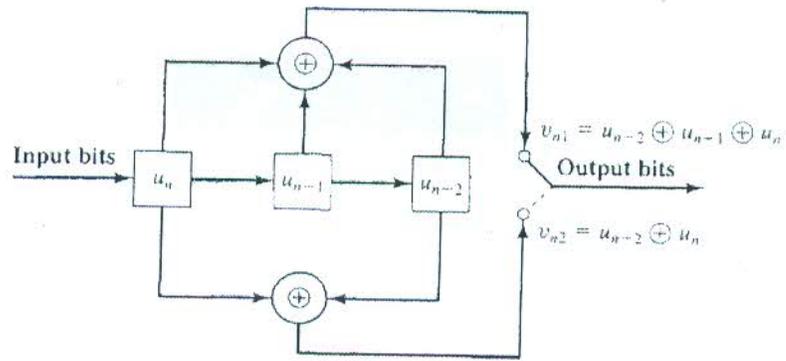
for convolutional code,

- useful when data are transmitted and received in a more or less continuous stream
- A convolutional code is defined by three parameters: n , k , and K .
- An (n, k, K) code processes input data k bits at a time and produces an output of n bits for each incoming k bits.
- n and k are generally quite small numbers.
- convolutional codes have memory, which is characterized by the *constraint factor* K . In essence, the current n -bit output of an (n, k, K) code depends not only on the value of the current block of k input bits but also on the previous $K - 1$ blocks of k input bits.

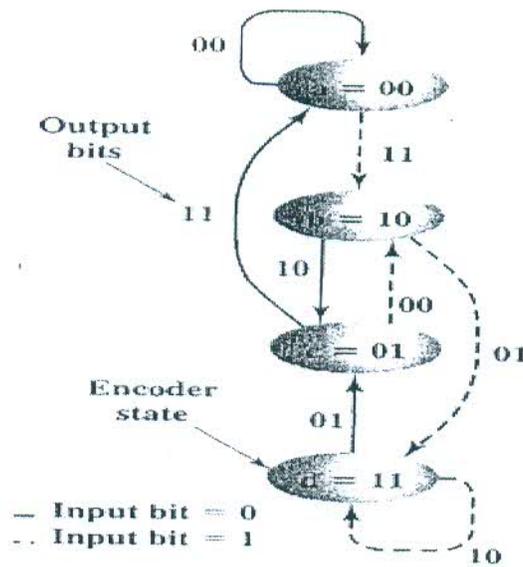
B) 7 marks

Encoder shift register

(12)



State diagram



c) 7 marks

to encode the following sequence 110010, we use the state diagram

Input	Previous state	Next state	output
0	00	00	00
1	00	10	11
0	10	01	10
0	01	00	11
1	00	10	11
1	10	11	01

Then the output sequence is

011111101100

from right to left

if we started from left

1101011110

[Q:5]

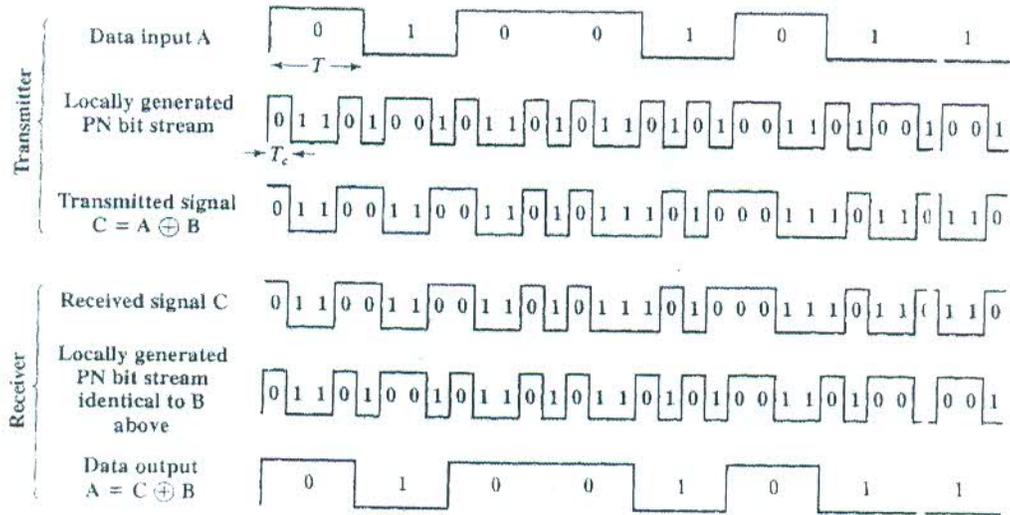
10 Marks

(13)

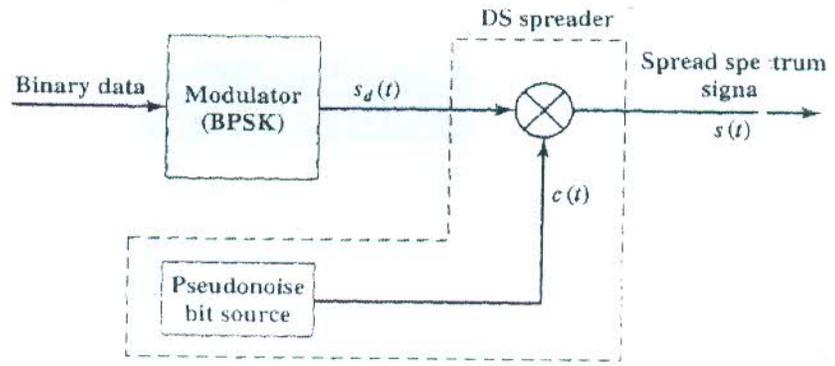
For direct sequence spread spectrum (DSSS), each bit in the original signal is represented by multiple bits in the transmitted signal, using a spreading code. The spreading code spreads the signal across a wider frequency band in direct proportion to the number of bits used. Therefore, a 10-bit spreading code spreads the signal across a frequency band that is 10 times greater than a 1-bit spreading code.

One technique for direct sequence spread spectrum is to combine the digital information stream with the spreading code bit stream using an exclusive-OR (XOR). The XOR obeys the following rules:

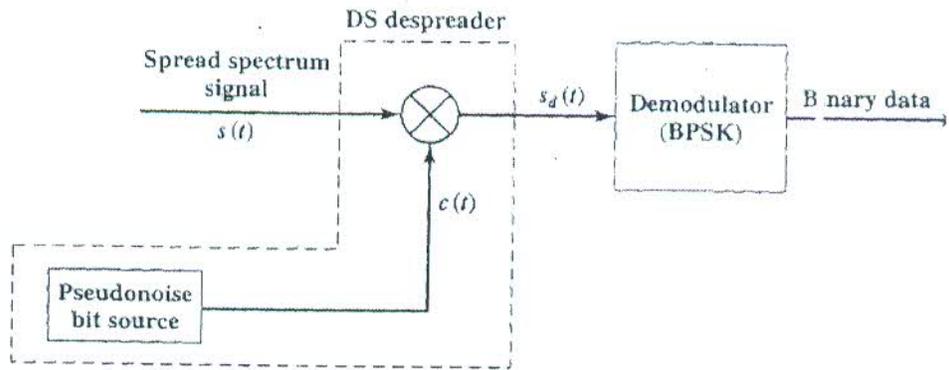
$$0 \oplus 0 = 0 \quad 0 \oplus 1 = 1 \quad 1 \oplus 0 = 1 \quad 1 \oplus 1 = 0$$



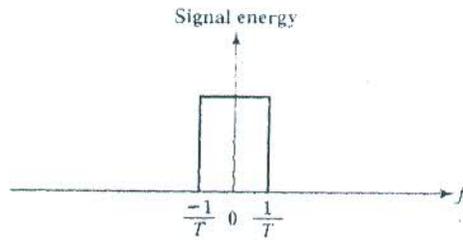
(19)



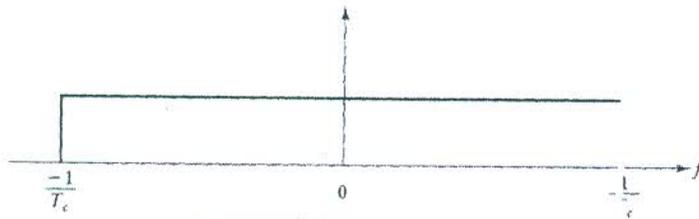
(a) Transmitter



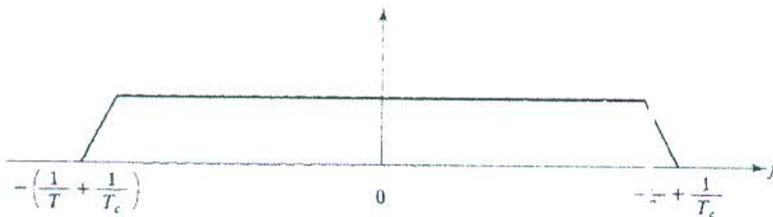
(b) Receiver



(a) Spectrum of data signal



(b) Spectrum of pseudonoise signal



(c) Spectrum of combined signal